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CITATION:

Tsuboi, Akito. Trees and Branching Axioms (Model Theory of Fields and its Applications).  
数理解析研究所講究録 2012, 1794: 50-54

ISSUE DATE:

2012-05

URL:

<http://hdl.handle.net/2433/172870>

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# Trees and Branching Axioms

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## 1 Introduction

First we recall the definition of trees. An ordered set  $O = (O, <)$  is called a tree if, for any  $a \in I$ , the initial segment  $O_a = \{b \in O : b < a\}$  is linearly ordered. A mapping  $\sigma : O \rightarrow O'$ , where  $O$  and  $O'$  are trees, is called a tree embedding if  $\sigma$  preserves  $<$ -structure, i.e.  $\eta < \nu$  if and only if  $\sigma(\eta) <' \sigma(\nu)$ . We are mainly interested in trees of the form  $\alpha^{<\beta}$ , where  $\alpha$  and  $\beta$  are ordinals and its order is  $<_{ini}$ :  $\eta <_{ini} \nu \iff \eta$  is a proper initial segment of  $\nu$ . The lexicographic order on  $\alpha^{<\beta}$  is denoted by  $<_{lex}$ . The meet operator  $\cap$  is a binary function that gives the greatest common lower bound.

We introduce the following notations:

- $A \simeq_{l.i.} B$  for expressing that  $A$  and  $B$  have the same  $\{<_{lex}, <_{ini}\}$ -atomic type.
- $A \simeq_{l.i.c.} B$  for expressing that  $A$  and  $B$  have the same  $\{<_{lex}, <_{ini}, \cap\}$ -atomic type.

Now let  $M$  be an  $L$ -structure. We consider a set  $A \subset M$  whose elements are indexed by a tree. So  $A$  has the form  $A = (a_\eta)_{\eta \in O}$ , where  $O$  is a tree. Such an indexed set is also called a tree. We introduce the notion of indiscernibility for such a tree  $A$ .

- $A$  is  $l.i$ -indiscernible if whenever  $X \simeq_{l.i.} Y$  then  $\text{tp}_L(a_X) = \text{tp}_L(a_Y)$ , where  $a_X = (a_\eta)_{\eta \in X}$ .
- $A$  is  $l.i.c$ -indiscernible if whenever  $X \simeq_{l.i.c.} Y$  then  $\text{tp}_L(a_X) = \text{tp}_L(a_Y)$ .

In this short note, we seek to find sufficient conditions for  $\Gamma(x_\eta)_{\eta \in O}$  to be realized by an indiscernible tree.

## 2 Indiscernible Trees

Throughout, let  $\sigma^* : \omega^{<\omega} \rightarrow \omega^{<\omega}$  be the mapping defined by

$$\langle m_0, \dots, m_{n-1} \rangle \mapsto \langle 0, m_0, \dots, 0, m_{n-1} \rangle.$$

This  $\sigma^*$  preserves  $<_{ini}$ , hence it is a tree embedding.  $<_{lex}$  is also preserved by  $\sigma^*$ .

**Remark 1** Let  $\eta, \nu$  be two  $<_{ini}$ -incomparable elements. Then  $\sigma^*(\eta \cap \nu)$  is a proper initial segment of  $\sigma^*(\eta) \cap \sigma^*(\nu)$ . So,  $A$  and  $\sigma^*A$  do not have the same *l.i.c.*-atomic type, unless  $A$  is linearly ordered.

**Definition 2** Let  $A \subset \omega^{<\omega}$  be a finite set. We say that  $A$  is a broom set if there are  $\eta_0, \dots, \eta_{n-1}$  such that

1.  $\eta_i \cap \eta_j = \eta_{i'} \cap \eta_{j'}$  for any  $i < j < n$  and  $i' < j' < n$ ,
2.  $A \subset \bigcup_{i < n} \{\eta_i | j : j \in \omega\}$ .

**Lemma 3** Let  $A, B \subset \omega^{<\omega}$ .

1. Suppose that  $A$  and  $B$  be broom sets. Then  $A \simeq_{l.i.} B \Rightarrow \sigma^*A \simeq_{l.i.c} \sigma^*B$ .
2. Suppose  $AC \simeq_{l.i.} BC$ , where  $A$  and  $B$  are broom sets. Suppose that for any incomparable  $\eta_1, \eta_2 \in A$  and any  $\eta \in C$ ,  $\eta_1 \cap \eta <_{ini} \eta_2 \cap \eta$ . Then  $\sigma^*(AC) \simeq_{l.i.c} \sigma^*(BC)$ .
3.  $A \simeq_{l.i.c} B \Rightarrow \sigma^*A \simeq_{l.i.c} \sigma^*B$ .

*Proof:* 2. We consider the most typical case, where  $A = \{\eta_1, \eta_2, \eta_3, \nu\}$ ,  $C = \{\eta\}$ ,  $\nu <_{ini} \eta_i$  ( $i = 1, 2, 3$ ),  $\nu <_{ini} \eta$  and  $\eta_1 \cap \eta_2 = \eta_2 \cap \eta_3 = \eta_3 \cap \eta_1$ . The *l.i.*-atomic type of  $\sigma^*(A)$  is determined by this data. Moreover, we have  $\sigma^*(\nu) <_{ini} \sigma^*(\eta_i) \cap \sigma^*(\eta_j)$  for any  $i < j$ , and  $\sigma^*(\nu) <_{ini} \sigma^*(\eta_i) \cap \sigma^*(\eta)$ . So the *l.i.c.*-atomic type of  $\sigma^*(A)$  is also determined. This argument proves  $A \simeq_{l.i.} B \Rightarrow \sigma^*A \simeq_{l.i.c} \sigma^*B$ .

3. Easy by the remark above.

Now we prepare the variables  $x_\eta$ , where  $\eta$  is a member of some fixed tree  $O$ . Usually, we are interested in the case  $O = \omega^{<\omega}$ . Let  $\Gamma((x_\eta)_{\eta \in \omega^{<\omega}})$  be a set of  $L$ -formulas with free variables from  $x_\eta$ 's.

**Definition 4** We say that  $\Gamma((x_\eta)_{\eta \in \omega^{<\omega}})$  has the subtree property if whenever  $I = (a_\eta)_{\eta \in \omega^{<\omega}}$  realizes  $\Gamma((x_\eta)_{\eta \in \omega^{<\omega}})$  and  $\sigma : \omega^{<\omega} \rightarrow \omega^{<\omega}$  is a tree embedding preserving *l.i.c.*-structure then  $I_\sigma = (a_{\sigma(\eta)})_{\eta \in \omega^{<\omega}}$  realizes  $\Gamma((x_\eta)_{\eta \in \omega^{<\omega}})$ .

**Lemma 5** Let  $\Gamma((x_\eta)_{\eta \in \omega^{<\omega}})$  be a consistent set having the subsequence property. Let  $\lambda$  be an infinite cardinal. Then there is a set  $J = (a_\eta)_{\eta \in \lambda^{<\omega}}$  such that for any  $\{<_{lex}, <_{ini}, <_{len}, P_n\}$ -embedding  $\sigma : \omega^{<\omega} \rightarrow \lambda^{<\omega}$  the set  $J_\sigma = (a_{\sigma(\eta)})_{\eta \in \omega^{<\omega}}$  realizes  $\Gamma((x_\eta)_{\eta \in \omega^{<\omega}})$ .

*Proof:* For  $A, B \subset \lambda^{<\omega}$ , we write  $A \simeq^+ B$  if  $A$  and  $B$  have the same atomic type in the language  $L_{l.i.c.l.} \cup \{P_n\}_{n \in \omega}$ . We prepare new variables  $x_\eta$  ( $\eta \in \lambda^{<\omega} \setminus \omega^{<\omega}$ ). Let  $\Gamma^*((x_\eta)_{\eta \in \lambda^{<\omega}})$  be the set obtained from  $\Gamma((x_\eta)_{\eta \in \omega^{<\omega}})$  by adding all formulas  $\varphi(x_A)$  with  $A \subset \lambda^{<\omega}$  such that  $\varphi(x_B) \in \Gamma((x_\eta)_{\eta \in \omega^{<\omega}})$  for some  $B \simeq^+ A$ . First we show

**Claim A**  $\Gamma^*$  is consistent.

Otherwise, there are  $\varphi_i(x_{A_i})$  and  $B_i$  ( $i < n$ ) such that

1.  $A_i \simeq^+ B_i$  and  $\varphi_i(x_{B_i}) \in \Gamma((x_\eta)_{\eta \in \omega^{<\omega}})$  ( $i < n$ ), and
2.  $\Gamma \vdash \bigvee_{i < n} \neg \varphi_i(x_{A_i})$ .

By compactness, there is a finite set  $\Gamma_0 \subset \Gamma$  such that  $\Gamma_0 \vdash \bigvee_{i < n} \neg \varphi_i(x_{A_i})$ . Hence, we can assume  $A_i$ 's are subsets of  $\omega^{<\omega}$ . Let  $N = \max\{\eta(n) : \eta \in \bigcup_i B_i, n \in \omega\}$  and let  $\sigma_N$  be the shift function mapping  $\eta = \langle \eta(0), \dots, \eta(n-1) \rangle$  to  $\langle \eta(0) + N, \dots, \eta(n-1) + N \rangle$ . Then, by the subtree property, we have

$$\Gamma((x_\eta)_{\eta \in \omega^{<\omega}}) \vdash \Gamma((x_{\sigma_N(\eta)})_{\eta \in \omega^{<\omega}}) \vdash \bigvee_{i < n} \neg \varphi_i(x_{\sigma_N(A_i)}).$$

From this, by replacing  $A_i$  with  $\sigma A_i$ , we can assume that  $A_i \subset (\omega \setminus N)^{<\omega}$ . Hence, for each  $i$ , there is a tree embedding  $\sigma_i$  that maps  $B_i$  to  $A_i$ . Choose a set  $(a_\eta)_{\eta \in \omega^{<\omega}}$  realizing  $\Gamma$ . By the property 2, there is  $i < n$  such that  $\neg \varphi(a_{A_i})$  holds. On the other hand, we have  $\varphi(x_{B_i}) \in \Gamma$  and  $\sigma_i(B_i) = A_i$ . Therefore, by the subtree property, we must have  $\varphi(a_{A_i})$ . A contradiction.

**Claim B** Let  $(a_\eta)_\eta$  be a realization of  $\Gamma^*$ . Then  $(a_\eta)_\eta$  has the desired condition.

**Lemma 6** Let  $\Gamma((x_\eta)_{\eta \in \omega^{<\omega}})$  be consistent and suppose that  $\Gamma$  has the subtree property. Then  $\Gamma$  is realized by an *l.i.c.-indiscernible tree*.

*Proof:* By Theorem 2.6 of [2, AP], since the width of the tree can be made arbitrarily large, we may assume that the tree  $(a_\eta)_{\eta \in \omega^{<\omega}}$  is an indiscernible tree in Shelah's sense. So, by Ramsey's theorem, we can choose an indiscernible tree  $I = (a_\eta)_{\eta \in \omega^{<\omega}}$  satisfying  $\Gamma$  such that if  $A$  and  $B$  have the same atomic type in the language  $L_{l.i.c.l.} = L_{l.i.c.} \cup \{<_{len}\}$  then  $a_A$  and  $a_B$  have the same  $L$ -type, where  $\eta <_{len} \nu$  means that the length of  $\eta$  is less than that of  $\nu$ .

By compactness, we can assume that the index set of  $I$  is  $\omega^{<\kappa}$ , where  $\kappa$  is very large. By induction on  $n \in \omega$ , we show that there is an *l.i.*-preserving mapping  $\sigma_n$  from  $\omega^{<n}$  to  $I$  such that if  $\eta <_{lex} \nu$  then  $\sigma_n(\eta) <_{len} \sigma_n(\nu)$ .

Suppose we have defined  $\sigma_n$ . Since  $\kappa$  is sufficiently large, there is  $\kappa_0 < \kappa$  such that the lengths of  $\sigma_n(\eta)$  ( $\eta \in \text{dom}(\sigma_n)$ ) are all less than  $\kappa_0$ . Now we define  $\sigma_{n+1}$  by the equation

$$\sigma_{n+1}(\langle i \rangle^\wedge \eta) = \underbrace{\langle i, i, \dots \rangle}_{\kappa_0 \cdot i}^\wedge \sigma_n(\eta).$$

This definition implies that  $\kappa_0 \cdot i \leq \text{len}(\sigma_{n+1}(\langle i \rangle^\wedge \eta)) < \kappa_0 \cdot (i + 1)$ . So, in particular, we have  $\text{len}(\sigma_{n+1}(\langle i \rangle^\wedge \eta)) < \text{len}(\sigma_{n+1}(\langle i' \rangle^\wedge \eta'))$ , if  $i < i'$ . By induction on the length of  $\eta$ , we can prove:

**Claim A**  $\sigma_{n+1}(\eta^\wedge \nu) = \sigma_n(\eta)^\wedge \sigma_n(\nu)$ , if  $\eta, \nu \in \text{dom}(\sigma_n)$ .

So,  $\sigma_{n+1}$  preserves *l.i.c.*-structure of the tree. Now we show:

**Claim B**  $\eta <_{lex} \eta' \Rightarrow \sigma_{n+1}(\eta) <_{len} \sigma_{n+1}(\eta')$ .

For proving this claim, let  $\nu = \eta \cap \eta'$ . If  $\eta <_{len} \eta'$  (i.e.  $\nu = \eta$ ), then clearly we have  $\sigma_{n+1}(\eta) <_{len} \sigma_{n+1}(\eta')$ . So we can assume  $\text{len}(\nu) > 0$ ,  $\eta = \nu^\wedge \langle i \rangle^\wedge \eta_0$ ,  $\eta' = \nu^\wedge \langle i' \rangle^\wedge \eta'_0$ , and  $i < i'$ . By Claim A, using the induction hypothesis, we have

$$\begin{aligned} \text{len}(\sigma_{n+1}(\eta)) &= \text{len}(\sigma_n(\nu)) + \text{len}(\sigma_n(\langle i \rangle^\wedge \eta_0)) \\ &< \text{len}(\sigma_n(\nu)) + \text{len}(\sigma_n(\langle i' \rangle^\wedge \eta'_0)) \\ &= \text{len}(\sigma_{n+1}(\eta')). \end{aligned}$$

Thus Claim B was shown, and  $\sigma_{n+1}$  has the required property. We have shown the existence of  $\sigma_n$ 's for all  $n$ . We fix  $n$  and put  $b_\eta = a_{\sigma_n(\eta)}$ . We prove:

**Claim C** *Let  $A, B \subset \text{dom}(\sigma_n)$  satisfy  $A \simeq_{l.i.c.} B$ . Then  $\text{tp}(b_A) = \text{tp}(b_B)$ .*

By  $A \simeq_{l.i.c.} B$ , we have  $\sigma_n(A) \simeq_{l.i.c.} \sigma_n(B)$ . So, by Claim B, we have

$$\sigma_n(A) \underset{l.i.c.l.}{\simeq} \sigma_n(B).$$

By the *l.i.c.l.*-indiscernibility of  $I$ , we have  $\text{tp}(a_{\sigma_n(A)}) = \text{tp}(a_{\sigma_n(B)})$ . Hence, from the definition  $b_\eta = a_{\sigma_n(\eta)}$ , we conclude  $\text{tp}(b_A) = \text{tp}(b_B)$ .

Now, by compactness and Claim C, we have the existence of *l.i.c.*-indiscernible trees realizing  $\Gamma$ .

**Theorem 7** *Let  $I = (a_\eta)_{\eta \in \omega^{<\omega}}$  be an *l.i.c.*-indiscernible tree. Let  $\sigma^*$  be the mapping described before. Let  $J = (b_\eta)_\eta = \sigma^* I$ .*

1.  *$J$  is an *l.i.c.*-indiscernible tree.*
2.  *$J$  is *l.i.*-indiscernible for broom sets: Suppose  $AC \simeq_{l.i.} BC$ , where  $A$  and  $B$  are broom sets. Suppose that for any incomparable  $\eta_1, \eta_2 \in A$  and any  $\nu \in C$ ,  $\eta_1 \cap \nu <_{ini} \eta_2 \cap \nu$ . Then  $\text{tp}((b_\eta)_{\eta \in AC}) = \text{tp}((b_\eta)_{\eta \in BC})$ .*

*Proof:* 1. Assume  $A \simeq_{l.i.c.} B$ . Then, by Lemma 3,  $\sigma^* A \simeq_{l.i.c.} \sigma^* B$ . By the tree indiscernibility, we have  $\text{tp}((a_\eta)_{\eta \in \sigma^* A}) = \text{tp}((a_\eta)_{\eta \in \sigma^* B})$ . The last equation is equivalent to

$$\text{tp}((a_{\sigma^*(\eta)})_{\eta \in A}) = \text{tp}((a_{\sigma^*(\eta)})_{\eta \in B}).$$

2. Clear by Lemma 3.

## References

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- [2] Saharon Shelah, *Classification Theory and the Number of Non-Isomorphic Models*, North-Holland, 1990